## Princeton University Department of Operations Research and Financial Engineering

## ORF 307 Optimization

## Practice Midterm 1

Closed book. No computers. Calculators allowed (but not needed).

You are permitted to use a one-page two-sided cheat sheet.

Please return the exam questions and your cheat sheet with your exam booklet.

(1) (30 pts.) Consider the following problem:

- (a) Write the problem in dictionary form.
- (b) Write down the values of all variables associated with this dictionary.
- (c) Write down the dual problem as a minimization problem.
- (d) Write the dual problem in dictionary form.
- (e) Write down the values of all variables associated with this dual dictionary.
- (f) Check whether the primal values from (b) are complementary to the dual values from (e).

(2) (14 pts.) Consider the following four dictionaries:

(a) 
$$\frac{\zeta = -2 - 2x_1 - 3w_1 - w_2}{x_2 = 2}$$

$$x_3 = 5 + 2x_1 - w_2$$
(b) 
$$\frac{\zeta = 3 + w_1 + 2x_2 - 2w_2 - x_3}{x_1 = 4}$$

$$x_4 = 3 - 4w_1 - x_3$$

For each of the above dictionaries, identify which of the following statements are true by placing check marks in the following table:

	(a)	(b)	(c)	(d)
The associated primal solution is feasible				
The associated primal solution is optimal				
The dictionary is primal degenerate				
The dictionary shows the primal problem is unbounded				
Variable $x_3$ is basic				
The associated dual solution is dual feasible				
The dictionary shows the primal problem is infeasible				

Be sure to return this page with your exam booklet.

(3) (18 pts.) Mark the following statements as true (T) or false (F):

	T/F
If the current dictionary is degenerate, then the next pivot will not change the objective function.	
Consider the version of the simplex method in which an artificial variable $x_0$ is introduced in phase-I. The phase-I problem can be unbounded.	
After a degenerate pivot, the dictionary must still be degenerate.	
An optimal solution to a linear programming problem must be a vertex of the set of feasible solutions.	
No problem with an unbounded feasible region has an optimal solution.	
Consider a dictionary. The primal and dual solutions associated with this dictionary must satisfy the complementarity relations: $x_j z_j = 0$ for all $j$ and $w_i y_i = 0$ for all $i$ .	
It is possible to move from one vertex of the feasible set to any other vertex in a single pivot.	
A dual pivot never decreases the value of the primal objective function.	
A problem has an optimal solution if and only if it is both primal and dual feasible.	

**NOTE:** SAT-style grading scheme: 2 points for a correct answer, 0 points for an incorrect answer, and 1 point for no answer.

(4) (7 pts.) The CEO of a large corporation needs the solution to the following linear programming problem:

maximize 
$$-2x_1 - 2x_2 - 6x_3$$
  
subject to  $x_1 - 2x_2 - 2x_3 \le -15$   
 $-2x_1 + x_2 - x_3 \le -30$   
 $x_1, x_2, x_3 \ge 0$ .

Because the numbers represent millions of dollars, an error would be very costly. Hence, he's independently delegated the task of computing the solution to two subordinates: Alice and Brad. Alice came back with the following values for the primal and, by the way, for the associated dual values:

$$x_1^* = 25, \ x_2^* = 20, \ x_3^* = 0$$
 and  $y_1^* = 2, \ y_2^* = 2$ 

Brad provided the following values, which are different from Alice's:

$$x_1^* = 9, \ x_2^* = 0, \ x_3^* = 12$$

Since the two solutions disagree, you're being asked to determine who, if anyone, is correct. So, who is? Explain.

(5) (7 pts.) Given n girls and m boys, you, the matchmaker, are to determine how to pairwise match the girls with the boys so as to maximize some overall desirability index. This problem can be set up as follows. Let  $x_{gb}$  be a decision variable that is 1 if girl g gets matched to boy g and is zero otherwise. Suppose that there is a desirability index g associated with matching girl g with boy g. The problem is to maximize the overall desirability of all of the matches subject to the constraints that each girl gets matched with one boy and each boy gets matched with at most one girl. Formulate this problem as a linear programming problem. Notes: (1) don't worry about specific data, just describe the problem; (2) you can write it either in AMPL or just in common mathematical notation; (3) don't worry that the solution might involve fractions.

(6) (8 pts.) The solution associated with the following dictionary is optimal:

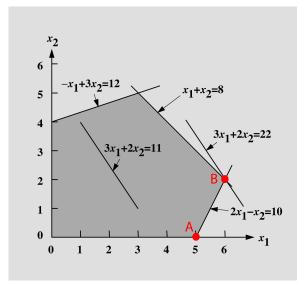
$$\begin{array}{cccccccc} \zeta & = 5 & - 2w_2 \\ \hline x_1 & = 6 + 5w_1 + w_2 \\ x_2 & = 0 - w_1 - 7w_2 \end{array}$$

- (a) Write down the optimal values for all of the variables in this problem.
- (b) If there are other optimal solutions, find one. If there are none, say how you know this to be the case.

(7) (6 pts.) Consider the following linear programming problem in 2 variables:

$$\begin{array}{lll} \text{maximize} & 3x_1 \, + \, 2x_2 \\ \text{subject to} & -x_1 \, + \, 3x_2 \, \leq \, 12 \\ & x_1 \, + \, \, x_2 \, \leq \, 8 \\ & 2x_1 \, - \, \, x_2 \, \leq \, 10 \\ & x_1, \, x_2 \, \geq \, \, 0. \end{array}$$

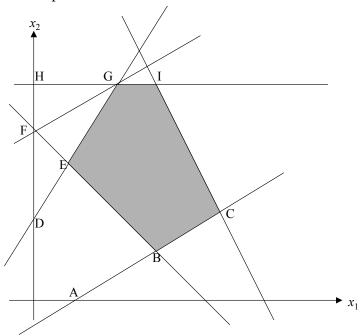
Here is a graph of the feasible region shown together with two level sets of the objective function:



Using the primal simplex method, one can pivot from the vertex labeled A to the vertex labeled B. Consider this pivot.

- (a) What is the entering variable?
- (b) What is the leaving variable? Explain.

(8) (5 pts.) Consider a problem whose constraints are as shown here:



Suppose that the objective is to minimize  $x_1$ . Clearly, the optimal solution is at vertex E. It is not possible for the dual simplex method starting at vertex H to visit vertices G then E? Explain why.

(9) (5 pts.) Using the largest coefficient rule to choose the entering variable in the following primal feasible dictionary, compute the dictionary after **one iteration** of the simplex method.